

## CLAIMS

1. Storage method of a plurality of data blocks in a digital re-writable memory of semiconductors controlled by a memory manager and characterized by the following steps that consist of:

- randomly determining an available area,
- storing the data block in the area chosen in this way.

2. Method according to claim 1, characterized in that it comprises a previous exploration step of the memory made by the memory manager, said exploration determining the available areas.

3. Method according to claim 2, characterized in that the result obtained after exploring the memory constitutes a list with the addresses of the available areas stored temporarily in a second memory, an address is then chosen randomly from said list and then the data block is stored in the area of the memory indicated by this address.

4. Method according to claim 2, characterized in that the exploration of the memory determines the number of available areas, a number between 1 and the number of areas found is randomly chosen and used to designate the area where the block must be stored.

5. Method according to claim 1, characterized in that a number N between 1 and the maximum number of available areas possible is randomly determined, the memory manager sequentially searches the N<sup>th</sup> available area and, if the end of the memory is reached before finding said area, the memory manager restarts the search from the beginning of the memory until it reaches the N<sup>th</sup> available area.

6. Method according to claims 1 to 5, characterized in that the blocks are of variable length, the storage of a block in the memory being made in an available area of a length equal or longer than the block length.

7. Method according to claims 1 to 5, characterized in that the data blocks are all the same length, the memory available areas having a length equal or longer than a multiple of the length of the blocks.
8. Method according to claims 1 to 6, characterized in that it includes a previous determination step of the usual length  $m$  of the blocks to be memorized, the blocks  $B_n$  of a shorter length than said usual value are stored at a pitch  $m$  from the preceding block in order to leave a free space equal to the difference between the usual length  $m$  and the block  $B_n$  length, the blocks  $B_n$  of equal length or longer than the current length  $m$  being stored immediately after the preceding block.
9. Method according to claim 1, characterized in that the memory is with direct access to the data through a pointer table, said pointers being chosen randomly before storage of the data blocks in the memory.
10. Method according to claim 9, characterized in that the pointer table is stored in a secured memory different from the main memory, each pointer being associated to an identifier of the data block.
11. Method according to claim 10, characterized in that said different secured memory is located in a removable security module such as a smart card.